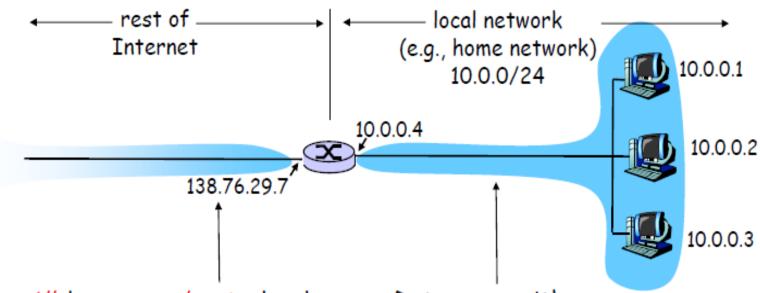
جامعة حماة الكلية التطبيقية قسم تقنيات حاسوب السنة الرابعة إدارة شبكات /٢/

م. إناس عدي المحاضرة السادسة



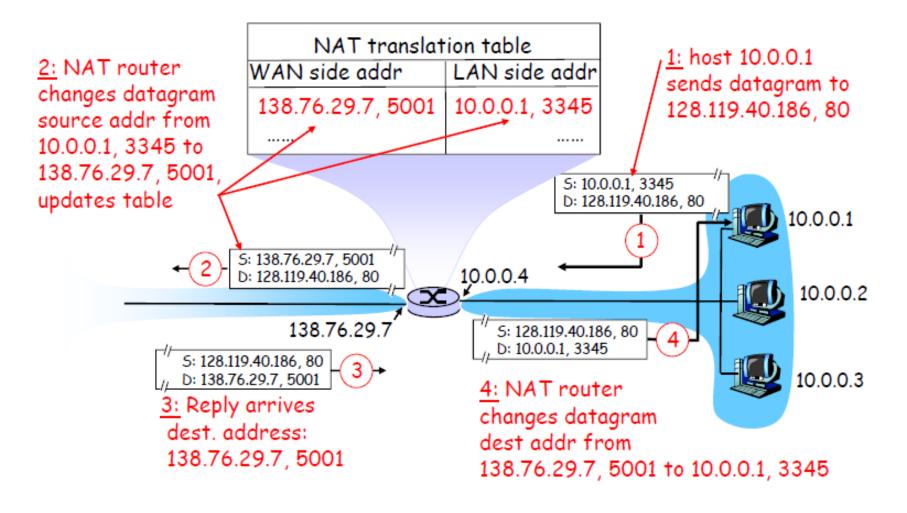
All datagrams leaving local network have same single source NAT IP address: 138.76.29.7, different source port numbers

Datagrams with source or destination in this network have 10.0.0/24 address for source, destination (as usual)

- Motivation: local network uses just one IP address as far as outside world is concerned:
 - range of addresses not needed from ISP: just one IP address for all devices
 - can change addresses of devices in local network without notifying outside world
 - can change ISP without changing addresses of devices in local network
 - devices inside local net not explicitly addressable, visible by outside world (a security plus).

Implementation: NAT router must:

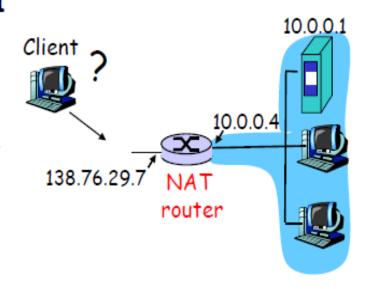
- outgoing datagrams: replace (source IP address, port #) of every outgoing datagram to (NAT IP address, new port #)
 - ... remote clients/servers will respond using (NAT IP address, new port #) as destination addr.
- remember (in NAT translation table) every (source IP address, port #) to (NAT IP address, new port #) translation pair
- incoming datagrams: replace (NAT IP address, new port #) in dest fields of every incoming datagram with corresponding (source IP address, port #) stored in NAT table



- □ 16-bit port-number field:
 - 60,000 simultaneous connections with a single LAN-side address!
- NAT is controversial:
 - o routers should only process up to layer 3
 - violates end-to-end argument
 - NAT possibility must be taken into account by app designers, eg, P2P applications
 - address shortage should instead be solved by IPv6

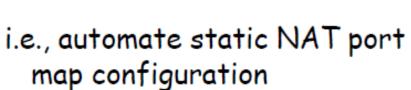
NAT traversal problem

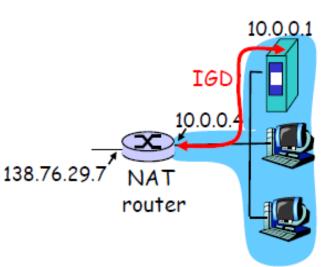
- client wants to connect to server with address 10.0.0.1
 - server address 10.0.0.1 local to LAN (client can't use it as destination addr)
 - only one externally visible NATted address: 138.76.29.7
- solution 1: statically configure NAT to forward incoming connection requests at given port to server
 - e.g., (123.76.29.7, port 2500) always forwarded to 10.0.0.1 port 25000



NAT traversal problem

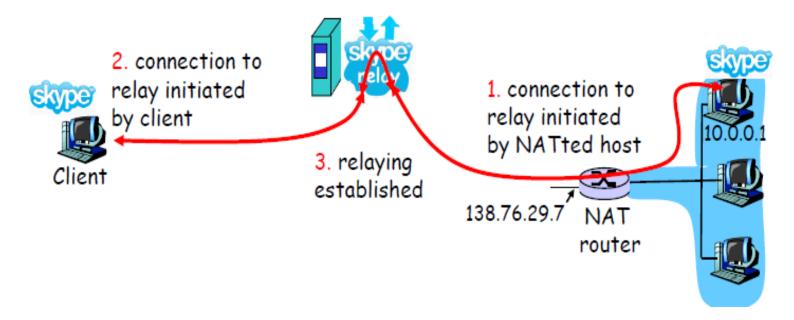
- solution 2: Universal Plug and Play (UPnP) Internet Gateway Device (IGD) Protocol. Allows NATted host to:
 - learn public IP address (138.76.29.7)
 - add/remove port mappings (with lease times)





NAT traversal problem

- solution 3: relaying (used in Skype)
 - NATed client establishes connection to relay
 - External client connects to relay
 - o relay bridges packets between to connections



ICMP: Internet Control Message Protocol

- used by hosts & routers to communicate network-level information
 - error reporting: unreachable host, network, port, protocol
 - echo request/reply (used by ping)
- □ network-layer "above" IP:
 - ICMP msgs carried in IP datagrams
- □ ICMP message: type, code plus first 8 bytes of IP datagram causing error

Туре	Code	description
0	0	echo reply (ping)
3	0	dest. network unreachable
3	1	dest host unreachable
3	2	dest protocol unreachable
3	3	dest port unreachable
3	6	dest network unknown
3	7	dest host unknown
4	0	source quench (congestion
		control - not used)
8	0	echo request (ping)
9	0	route advertisement
10	0	router discovery
11	0	TTL expired
12	0	bad IP header

Traceroute and ICMP

- Source sends series of UDP segments to dest
 - First has TTL =1
 - Second has TTL=2, etc.
 - Unlikely port number
- When nth datagram arrives to nth router:
 - Router discards datagram
 - And sends to source an ICMP message (type 11, code 0)
 - Message includes name of router& IP address

- When ICMP message arrives, source calculates RTT
- Traceroute does this 3 times

Stopping criterion

- UDP segment eventually arrives at destination host
- Destination returns ICMP "host unreachable" packet (type 3, code 3)
- When source gets this ICMP, stops.

<u>IPv6</u>

- □ Initial motivation: 32-bit address space soon to be completely allocated.
- Additional motivation:
 - header format helps speed processing/forwarding
 - header changes to facilitate QoS

IPv6 datagram format:

- o fixed-length 40 byte header
- no fragmentation allowed

IPv6 Header (Cont)

Priority: identify priority among datagrams in flow Flow Label: identify datagrams in same "flow."

(concept of "flow" not well defined).

Next header: identify upper layer protocol for data

ver pri	flow label						
payload	l len	next hdr	hop limit				
source address (128 bits)							
destination address							
(128 bits)							
data							
22 bits							

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Other Changes from IPv4

- Checksum: removed entirely to reduce processing time at each hop
- Options: allowed, but outside of header, indicated by "Next Header" field
- □ ICMPv6: new version of ICMP
 - o additional message types, e.g. "Packet Too Big"
 - multicast group management functions

Transition From IPv4 To IPv6

- Not all routers can be upgraded simultaneous
 - ono "flag days"
 - O How will the network operate with mixed IPv4 and IPv6 routers?
- Tunneling: IPv6 carried as payload in IPv4 datagram among IPv4 routers

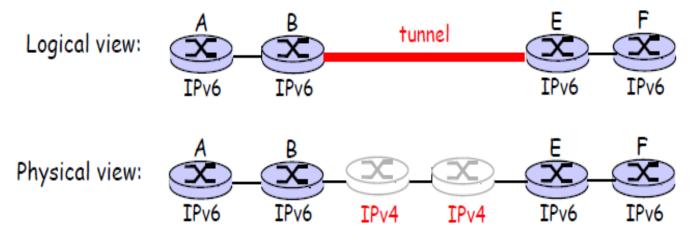
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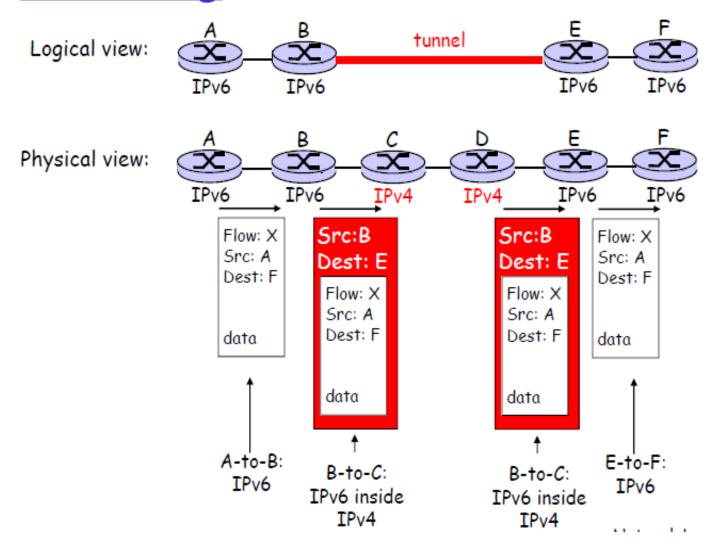
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Tunneling



Tunneling



The End